## Alternator Coins

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## Original Coin Problem



You are given $N$ coins that look identical, but one of them is fake and is lighter than other coins. You have a balance scale that you can use to help find the fake coin. What is the smallest number of weighings that guarantees finding the fake coin?

## Example

This is a coin problem that first appeared in 1945. Since then, there were many generalizations of this puzzle.

Try this problem: What is the smallest number of weighings that guarantees finding the fake coin from a group of eight coins?

## Answer

Answer: 2

## YAY!

## Original Coin Problem - Solution to All Cases

For a case with $N$ coins, the number of weighing will be $\left\lceil\log _{3} N\right\rceil$.

## Alternator Coin Idea

Alternator Coin: A coin that starts out randomly: fake or real, and then after each weighing that it participates in, it switches state.


## Alternator Coin — States

f-state - The alternator coin will act as a fake coin in its next weighing.
r-state - The alternator coin will act as a real coin in its next weighing.

## Solutions

$N$ is the total number of coins.
$f(N)$ - The smallest number of weighings to find the alternator if the alternator coin is currently in the f-state.
$r(N)$ - The smallest number of weighings to find the alternator if the alternator coin is currently in the r-state.
$a(N)$ - smallest number of weighings to find the alternator if the state of the alternator is unknown.


## Trivial Bounds

There are trivial lower and upper bounds:

- Lower bound: the alternator is worse than the fake coin.
- Upper bound: the alternator is better than two times the weighings needed for one fake coin.


## Example: 3 coins

What are $r(N)$ and $f(N)$ for 3 coins?


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$$
\begin{aligned}
& \mathrm{r}(N)=2 \\
& \mathrm{f}(N)=1
\end{aligned}
$$

## Jacobsthal Numbers

Sequence $J_{n}: 0,1,1,3,5,11,21,43 \ldots$ Can you guess the rule?

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$$
J_{n}=\left(2^{n}-(-1)^{n}\right) / 3
$$

## Observations

We made the following observations:

- The number of weighings necessary increases by one after the number of coins reaches the next Jacobsthal number.
- $f(N)$ is always equal to $r(N)-1$.
- $r(N)=a(N)$.


## Ideal Strategy

There are 11 coins, one of which is an alternator coin. How many weighings on a two pan balance will it take to find the alternator coin?


## Strategy: $f$-state

$$
\begin{aligned}
& 11=J_{5} \\
& J_{5}=J_{4}+2 J_{3}: 11=5+2 \cdot 3 .
\end{aligned}
$$

$$
\text { We compare } 3 \text { coins versus } 3 \text { coins. }
$$

- If unbalances: $r(3)=2$.
- If balances: $f(5)=2$.

Thus, $f(11)=3$.

## Strategy: r-state

- Even number of coins: put all of them on the scale: $r(2 k)=f(2 k)+1$.
- Odd number of coins: put one aside. Later if everything balances, then this is the alternator: $r(2 k+1)=f(2 k)+1$.


## Optimum I

- E: equal
- L: left pan is heavier
- R: right pan is heavier

Every unique string points to a different coin.

## Optimum II

Property: L or R must be followed by E.
The number of such strings of length $n$ is $J_{n+2}$ :

- Length 0 : one string: only empty string.
- Length 1: three strings, E,L,R.

The number of such strings, $s(n)$ :
$s(n)=s(n-1)+2 s(n-2)$,
For the $r$-state the string has to start with E , so the number of such strings of length $n$ is $J_{n+1}$.

## Results

## Theorem

For the $f$ state, the number of coins $N$ we can process in $w$ weighings is $J_{w+1}<N \leq J_{w+2}$. For the $r$ state, the number of coins $N$ we can process in $w$ weighings is $J_{w}<N \leq J_{w+1}$.

## Theorem

For the a-state and $r$-state, the number of coins $N$ we can process in $w$ weighings is $J_{w}<N \leq J_{w+1}$.

## Corollary

$$
a(N)=r(N)=f(N)+1 .
$$

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